- 1. Consider an *n*-node treap *T*. As in the lecture notes, we identify nodes in *T* by the ranks of their search keys; for example, 'node 5' means the node with the 5th smallest search key. Let i, j, k be integers such that $1 \le i \le j \le k \le n$.
 - (a) What is the *exact* probability that node *j* is a common ancestor of node *i* and node *k*?
 - (b) What is the *exact* expected length of the unique path in *T* from node *i* to node *k*?

Don't forget to prove that your answers are correct!

- 2. A *meldable priority queue* stores a set of keys from some totally-ordered universe (such as the integers) and supports the following operations:
 - MAKEQUEUE: Return a new priority queue containing the empty set.
 - FINDMIN(*Q*): Return the smallest element of *Q* (if any).
 - DELETEMIN(*Q*): Remove the smallest element in *Q* (if any).
 - INSERT(Q, x): Insert element x into Q, if it is not already there.
 - DECREASEKEY(Q, x, y): Replace an element $x \in Q$ with a smaller key y. (If y > x, the operation fails.) The input is a pointer directly to the node in Q containing x.
 - DELETE(*Q*, *x*): Delete the element *x* ∈ *Q*. The input is a pointer directly to the node in *Q* containing *x*.
 - MELD (Q_1, Q_2) : Return a new priority queue containing all the elements of Q_1 and Q_2 ; this operation destroys Q_1 and Q_2 .

A simple way to implement such a data structure is to use a heap-ordered binary tree, where each node stores a key, along with pointers to its parent and two children. MELD can be implemented using the following randomized algorithm:

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\begin{array}{c} \underline{\operatorname{MELD}(Q_1,Q_2):}\\ \text{if } Q_1 \text{ is empty return } Q_2\\ \text{if } Q_2 \text{ is empty return } Q_1\\ \text{if } key(Q_1) > key(Q_2)\\ \text{swap } Q_1 \leftrightarrow Q_2\\ \text{with probability } 1/2\\ left(Q_1) \leftarrow \operatorname{MELD}(left(Q_1),Q_2)\\ \text{else}\\ right(Q_1) \leftarrow \operatorname{MELD}(right(Q_1),Q_2)\\ \text{return } Q_1 \end{array}
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- (a) Prove that for *any* heap-ordered binary trees Q_1 and Q_2 (not just those constructed by the operations listed above), the expected running time of MELD (Q_1, Q_2) is $O(\log n)$, where $n = |Q_1| + |Q_2|$. [Hint: How long is a random root-to-leaf path in an *n*-node binary tree if each left/right choice is made with equal probability?]
- (b) Show that each of the other meldable priority queue operations can be implemented with at most one call to MELD and O(1) additional time. (This implies that every operation takes $O(\log n)$ expected time.)