1. A *vertex cover* of a graph is a subset *S* of the vertices such that every vertex *v* either belongs to *S* or has a neighbor in *S*. In other words, the vertices in *S* cover all the edges. Finding the minimum size of a vertex cover is *NP*-hard, but in trees it can be found using dynamic programming.

Given a tree *T* and non-negative weight w(v) for each vertex *v*, describe an algorithm computing the minimum weight of a vertex cover of *T*.

2. Suppose you are given an unparenthesized mathematical expression containing *n* numbers, where the only operators are + and -; for example:

$$1 + 3 - 2 - 5 + 1 - 6 + 7$$

You can change the value of the expression by adding parentheses in different positions. For example:

$$1+3-2-5+1-6+7 = -1$$
$$(1+3-(2-5))+(1-6)+7 = 9$$
$$(1+(3-2))-(5+1)-(6+7) = -17$$

Design an algorithm that, given a list of integers separated by + and - signs, determines the maximum possible value the expression can take by adding parentheses.

You can only insert parentheses immediately before and immediately after numbers; in particular, you are not allowed to insert implicit multiplication as in 1 + 3(-2)(-5) + 1 - 6 + 7 = 33.

3. Fix an arbitrary sequence  $c_1 < c_2 < \cdots < c_k$  of coin values, all in cents. We have an infinite number of coins of each denomination. Describe a dynamic programming algorithm to determine, given an arbitrary non-negative integer x, the least number of coins whose total value is x. For simplicity, you may assume that  $c_1 = 1$ .

## To think about later after learning "greedy algorithms":

- (a) Describe a greedy algorithm to make change consisting of quarters, dimes, nickels, and pennies. Prove that your algorithm yields an optimal solution.
- (b) Suppose that the available coins have the values  $c^0, c^1, ..., c^k$  for some integers c > 1 and  $k \ge 1$ . Show that the greedy algorithm always yields an optimal solution.
- (c) Describe a set of 4 coin values for which the greedy algorithm does *not* yield an optimal solution.