

# CS 373: Combinatorial Algorithms, Spring 1999

## Final Exam (May 7, 1999)

Name:	
Net ID:	Alias:

**This is a closed-book, closed-notes exam!**

If you brought anything with you besides writing instruments and your two  $8\frac{1}{2}'' \times 11''$  cheat sheets, please leave it at the front of the classroom.

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- Print your name, netid, and alias in the boxes above, and print your name at the top of every page.
  - **Answer six of the seven questions on the exam.** Each question is worth 10 points. If you answer every question, the one with the lowest score will be ignored. **1-unit graduate students must answer question #7.**
  - Please write your answers on the front of the exam pages. Use the backs of the pages as scratch paper. Let us know if you need more paper.
  - Read the entire exam before writing anything. Make sure you understand what the questions are asking. If you give a beautiful answer to the wrong question, you'll get no credit. If any question is unclear, please ask one of us for clarification.
  - Don't spend too much time on any single problem. If you get stuck, move on to something else and come back later.
  - Write *something* down for every problem. Don't panic and erase large chunks of work. Even if you think it's nonsense, it might be worth partial credit.
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#	Score	Grader
1		
2		
3		
4		
5		
6		
7		

## 1. Short Answer

sorting	induction	Master theorem	divide and conquer
randomized algorithm	amortization	brute force	hashing
binary search	depth-first search	splay tree	Fibonacci heap
convex hull	sweep line	minimum spanning tree	shortest paths
shortest path	adversary argument	NP-hard	reduction
string matching	evasive graph property	dynamic programming	$H_n$

Choose from the list above the best method for solving each of the following problems. We do *not* want complete solutions, just a short description of the proper solution technique! Each item is worth 1 point.

- Given a Champaign phone book, find your own phone number.
- Given a collection of  $n$  rectangles in the plane, determine whether any two intersect in  $O(n \log n)$  time.
- Given an undirected graph  $G$  and an integer  $k$ , determine if  $G$  has a complete subgraph with  $k$  edges.
- Given an undirected graph  $G$ , determine if  $G$  has a triangle — a complete subgraph with three vertices.
- Prove that any  $n$ -vertex graph with minimum degree at least  $n/2$  has a Hamiltonian cycle.
- Given a graph  $G$  and three distinguished vertices  $u$ ,  $v$ , and  $w$ , determine whether  $G$  contains a path from  $u$  to  $v$  that passes through  $w$ .
- Given a graph  $G$  and two distinguished vertices  $u$  and  $v$ , determine whether  $G$  contains a path from  $u$  to  $v$  that passes through at most 17 edges.
- Solve the recurrence  $T(n) = 5T(n/17) + O(n^{4/3})$ .
- Solve the recurrence  $T(n) = 1/n + T(n - 1)$ , where  $T(0) = 0$ .
- Given an array of  $n$  integers, find the integer that appears most frequently in the array.

(a) \_\_\_\_\_ (f) \_\_\_\_\_

(b) \_\_\_\_\_ (g) \_\_\_\_\_

(c) \_\_\_\_\_ (h) \_\_\_\_\_

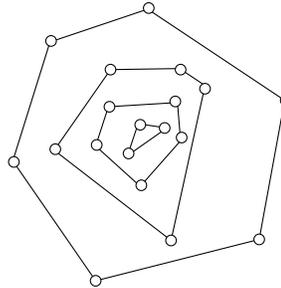
(d) \_\_\_\_\_ (i) \_\_\_\_\_

(e) \_\_\_\_\_ (j) \_\_\_\_\_

## 2. Convex Layers

Given a set  $Q$  of points in the plane, define the *convex layers* of  $Q$  inductively as follows: The first convex layer of  $Q$  is just the convex hull of  $Q$ . For all  $i > 1$ , the  $i$ th convex layer is the convex hull of  $Q$  after the vertices of the first  $i - 1$  layers have been removed.

Give an  $O(n^2)$ -time algorithm to find all convex layers of a given set of  $n$  points. [Partial credit for a correct slower algorithm; extra credit for a correct faster algorithm.]

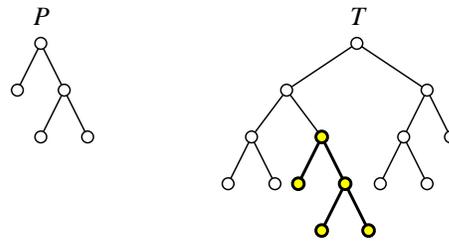


A set of points with four convex layers.

3. Suppose you are given an array of  $n$  numbers, sorted in increasing order.
- (a) **[3 pts]** Describe an  $O(n)$ -time algorithm for the following problem:  
Find two numbers from the list that add up to zero, or report that there is no such pair. In other words, find two numbers  $a$  and  $b$  such that  $a + b = 0$ .
- (b) **[7 pts]** Describe an  $O(n^2)$ -time algorithm for the following problem:  
Find *three* numbers from the list that add up to zero, or report that there is no such triple. In other words, find three numbers  $a$ ,  $b$ , and  $c$ , such that  $a + b + c = 0$ . [Hint: Use something similar to part (a) as a subroutine.]

## 4. Pattern Matching

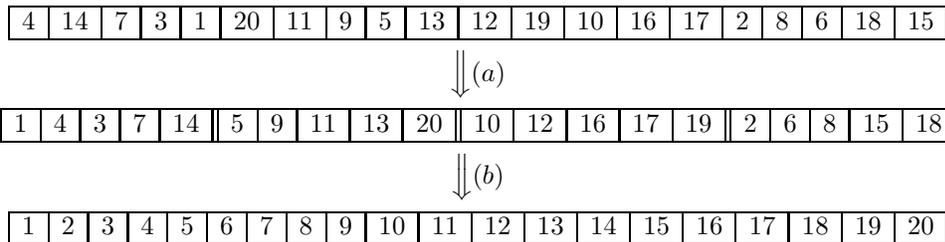
- (a) [4 pts] A *cyclic rotation* of a string is obtained by chopping off a prefix and gluing it at the end of the string. For example, ALGORITHM is a cyclic shift of RITHMALGO. Describe and analyze an algorithm that determines whether one string  $P[1..m]$  is a cyclic rotation of another string  $T[1..n]$ .
- (b) [6 pts] Describe and analyze an algorithm that decides, given any two binary trees  $P$  and  $T$ , whether  $P$  equals a subtree of  $T$ . [Hint: First transform both trees into strings.]



$P$  occurs exactly once as a subtree of  $T$ .

## 5. Two-stage Sorting

- (a) [1 pt] Suppose we are given an array  $A[1..n]$  of distinct integers. Describe an algorithm that splits  $A$  into  $n/k$  subarrays, each with  $k$  elements, such that the elements of each subarray  $A[(i-1)k+1..ik]$  are sorted. Your algorithm should run in  $O(n \log k)$  time.
- (b) [2 pts] Given an array  $A[1..n]$  that is already split into  $n/k$  sorted subarrays as in part (a), describe an algorithm that sorts the entire array in  $O(n \log(n/k))$  time.
- (c) [3 pts] Prove that your algorithm from part (a) is optimal.
- (d) [4 pts] Prove that your algorithm from part (b) is optimal.



**6. SAT Reduction**

Suppose you are have a black box that magically solves SAT (the formula satisfiability problem) in constant time. That is, given a boolean formula of variables and logical operators ( $\wedge, \vee, \neg$ ), the black box tells you, in constant time, whether or not the formula can be satisfied. Using this black box, design and analyze a **polynomial-time** algorithm that computes an assignment to the variables that satisfies the formula.

**7. Knapsack**

You're hiking through the woods when you come upon a treasure chest filled with objects. Each object has a different size, and each object has a price tag on it, giving its value. There is no correlation between an object's size and its value. You want to take back as valuable a subset of the objects as possible (in one trip), but also making sure that you will be able to carry it in your knapsack which has a limited size.

In other words, you have an integer capacity  $K$  and a target value  $V$ , and you want to decide whether there is a subset of the objects whose total size is *at most*  $K$  and whose total value is *at least*  $V$ .

- (a) **[5 pts]** Show that this problem is NP-hard. [Hint: Restate the problem more formally, then reduce from the NP-hard problem PARTITION: Given a set  $S$  of nonnegative integers, is there a partition of  $S$  into disjoint subsets  $A$  and  $B$  (where  $A \cup B = S$ ) whose sums are equal, *i.e.*,  $\sum_{a \in A} a = \sum_{b \in B} b$ .]
- (b) **[5 pts]** Describe and analyze a dynamic programming algorithm to solve the knapsack problem in  $O(nK)$  time. Prove your algorithm is correct.