

“CS 374” Fall 2014 ✧ Homework 11

Due Tuesday, December 9, 2014 at noon

1. Recall that w^R denotes the reversal of string w ; for example, $\text{TURING}^R = \text{GNIRUT}$. Prove that the following language is undecidable.

$$\text{REVACCEPT} := \{ \langle M \rangle \mid M \text{ accepts } \langle M \rangle^R \}$$

2. Let M be a Turing machine, let w be an arbitrary input string, and let s be an integer. We say that M **accepts w in space s** if, given w as input, M accesses only the first s cells on the tape and eventually accepts.

- (a) Prove that the following language is decidable:

$$\{ \langle M, w \rangle \mid M \text{ accepts } w \text{ in space } |w|^2 \}$$

- (b) Prove that the following language is undecidable:

$$\{ \langle M \rangle \mid M \text{ accepts at least one string } w \text{ in space } |w|^2 \}$$

3. **[Extra credit]** For each of the following languages, either prove that the language is decidable, or prove that the language is undecidable.

(a) $L_0 = \{ \langle M \rangle \mid \text{given any input string, } M \text{ eventually leaves its start state} \}$

(b) $L_1 = \{ \langle M \rangle \mid M \text{ decides } L_0 \}$

(c) $L_2 = \{ \langle M \rangle \mid M \text{ decides } L_1 \}$

(d) $L_3 = \{ \langle M \rangle \mid M \text{ decides } L_2 \}$

(e) $L_4 = \{ \langle M \rangle \mid M \text{ decides } L_3 \}$

CS 374 Fall 2014 ✧ Homework 11 Problem 1

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Prove that $\text{REVACCEPT} := \{\langle M \rangle \mid M \text{ accepts } \langle M \rangle^R\}$ is undecidable.

CS 374 Fall 2014 ✧ Homework 11 Problem 2

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- (a) Prove that $\{\langle M, w \rangle \mid M \text{ accepts } w \text{ in space } |w|^2\}$ is decidable.
- (b) Prove that $\{\langle M \rangle \mid M \text{ accepts at least one string } w \text{ in space } |w|^2\}$ is undecidable:
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CS 374 Fall 2014 ✧ Homework 11 Problem 3
[Extra credit]

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For each of the following languages, either prove that the language is decidable, or prove that the language is undecidable.

- (a) $L_0 = \{ \langle M \rangle \mid \text{given any input string, } M \text{ eventually leaves its start state} \}$
 - (b) $L_1 = \{ \langle M \rangle \mid M \text{ decides } L_0 \}$
 - (c) $L_2 = \{ \langle M \rangle \mid M \text{ decides } L_1 \}$
 - (d) $L_3 = \{ \langle M \rangle \mid M \text{ decides } L_2 \}$
 - (e) $L_4 = \{ \langle M \rangle \mid M \text{ decides } L_3 \}$
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