

CS 473 ✧ Spring 2017

☪ Homework 9 ☪

Due Wednesday, April 19, 2017 at 8pm

1. Suppose you are given an arbitrary directed graph $G = (V, E)$ with arbitrary edge weights $\ell: E \rightarrow \mathbb{R}$. Each edge in G is colored either red, white, or blue to indicate how you are permitted to modify its weight:

- You may increase, but not decrease, the length of any red edge.
- You may decrease, but not increase, the length of any blue edge.
- You may not change the length of any black edge.

The *cycle nullification* problem asks whether it is possible to modify the edge weights—subject to these color constraints—so that *every cycle in G has length 0*. Both the given weights and the new weights of the individual edges can be positive, negative, or zero. To keep the following problems simple, assume that G is strongly connected.

- (a) Describe a linear program that is feasible if and only if it is possible to make every cycle in G have length 0. [Hint: Pick an arbitrary vertex s , and let $\text{dist}(v)$ denote the length of every walk from s to v .]
 - (b) Construct the dual of the linear program from part (a). [Hint: Choose a convenient objective function for your primal LP.]
 - (c) Give a self-contained description of the combinatorial problem encoded by the dual linear program from part (b), and prove *directly* that it is equivalent to the original cycle nullification problem. Do not use the words “linear”, “program”, or “dual”. Yes, you have seen this problem before.
 - (d) Describe and analyze an algorithm to determine *in $O(EV)$ time* whether it is possible to make every cycle in G have length 0, using your dual formulation from part (c). Do not use the words “linear”, “program”, or “dual”.
2. *There is no problem 2.*