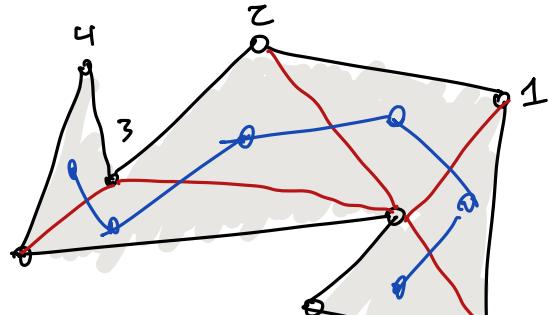
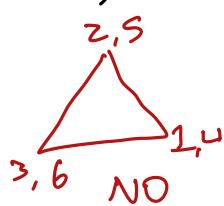
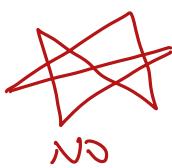


Polygon Triangulation

Input: Simple polygon in plane P



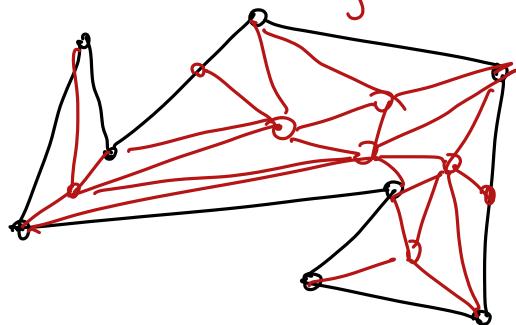
Specified as sequence of vertices in cyclic order



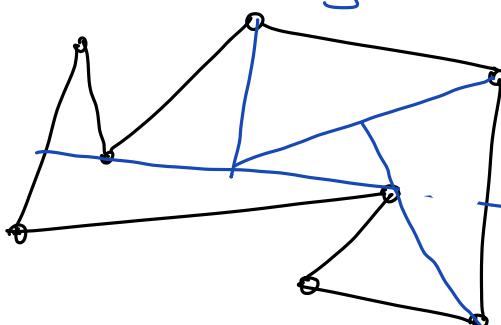
Compute a (frugal, interior) triangulation of P

↑
Jordan curve theorem (1906)

Not Frugal

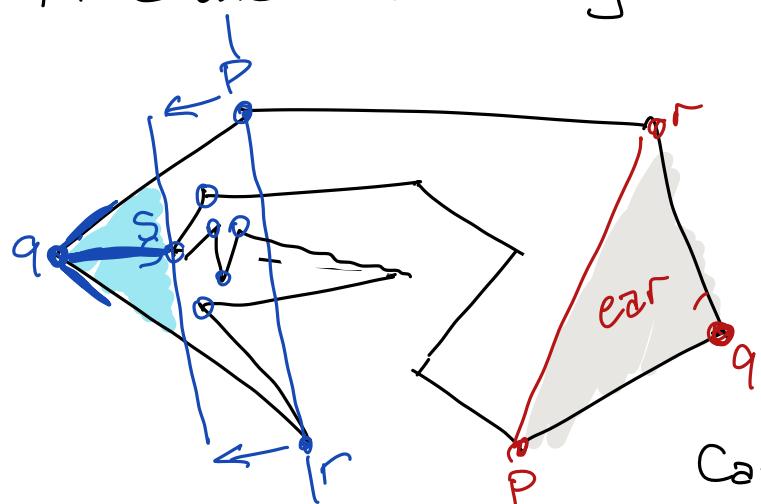


Not triangulation



→ Lennes (1911)
Dehn (1899) — unpublished

Find one interior diagonal



Lemma: Every simple polygon with > 3 vertices has an interior diagonal

Proof:

Pick an extreme vertex q
 $p, r \leftarrow \text{pred, succ of } q$

Case ① pr is interior diagonal ✓

case ②: pr is not an interior diagonal

P has vertices inside Δpr

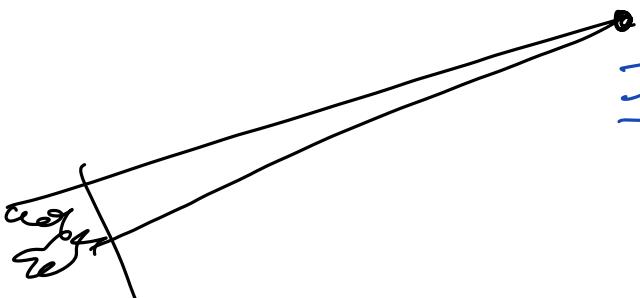
connect q to one of these vertices

qs is int. diagonal! $s = \text{interior, vertex Furthest from } pr$

Algo: we can find a diagonal in $O(n)$ time

\Rightarrow we can compute a triangulation in $O(n^2)$ time

$$T(n) = T(k) + T(n+2-k) + O(n)$$
$$= O(n^2) \text{ worst case}$$



Interesting problem:

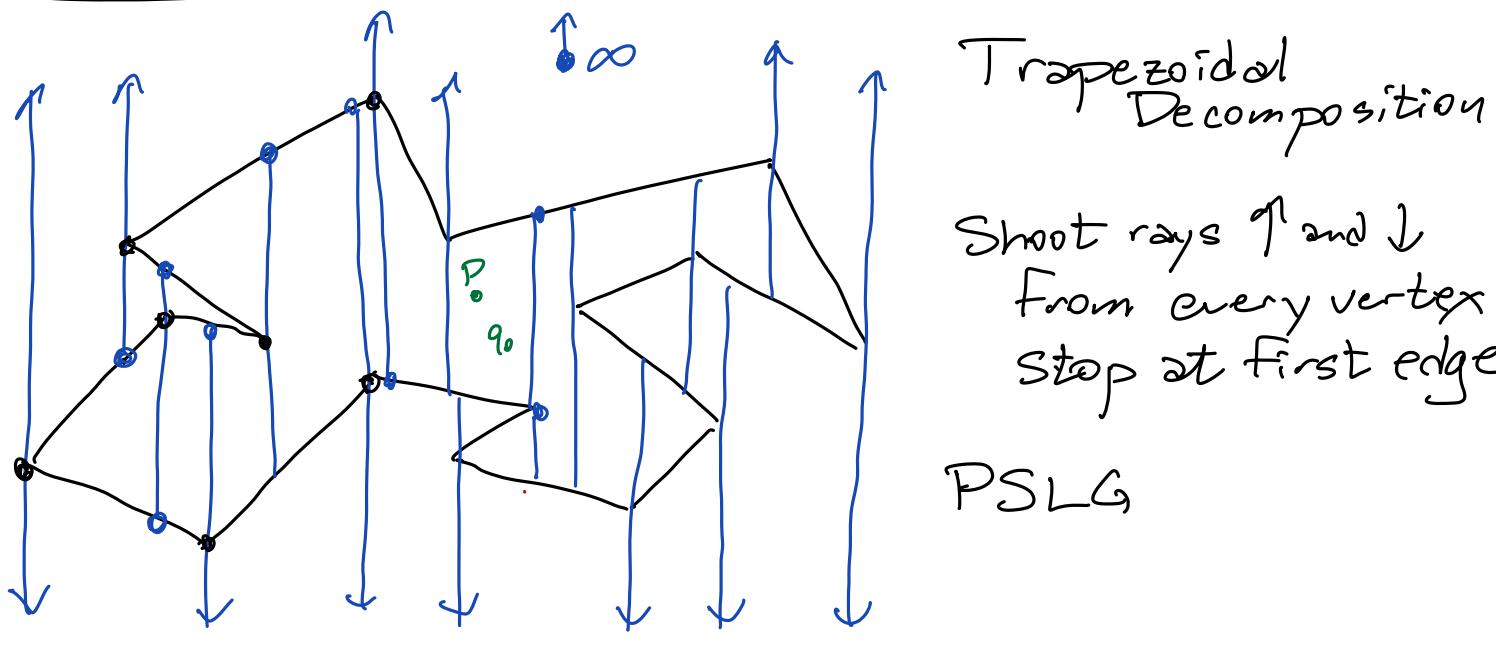
Find 2 diagonal splits vertices $\frac{n}{3} + \frac{2n}{3}$ or better.

$O(n \log n)$ today \leftarrow sweep line

$O(n \log^2 n)$ later

$O(n)$ Chazelle '90

Amato Goodrich Ramos OS?



Events = vertices

$\nearrow -\infty$

At each event in \rightarrow order

find succ + pred

build walls up and down

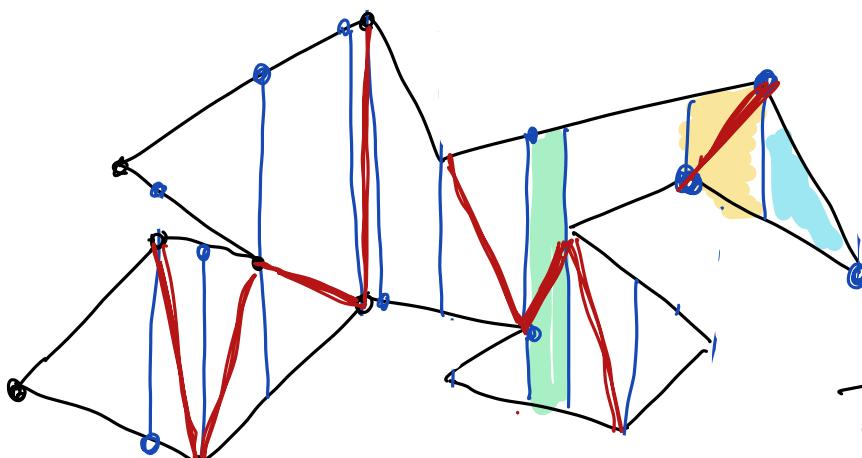
add $O(1)$ verts and edges to trap-decomp graph.

We can build trap-decomp. in $O(n \log n)$ time using a sweep algo.

Define $p \sim q$ ("equivalent") if same edges above and below.

Trapezoid = equivalence class!

How to triangulate



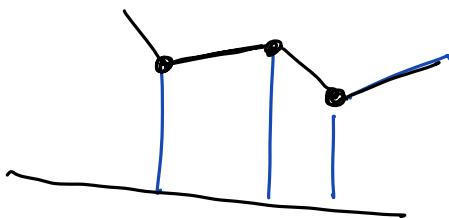
- ① Build trap decompo
- ② discard exterior walls

Every trap has
2 vertex on L wall
vertex on R wall

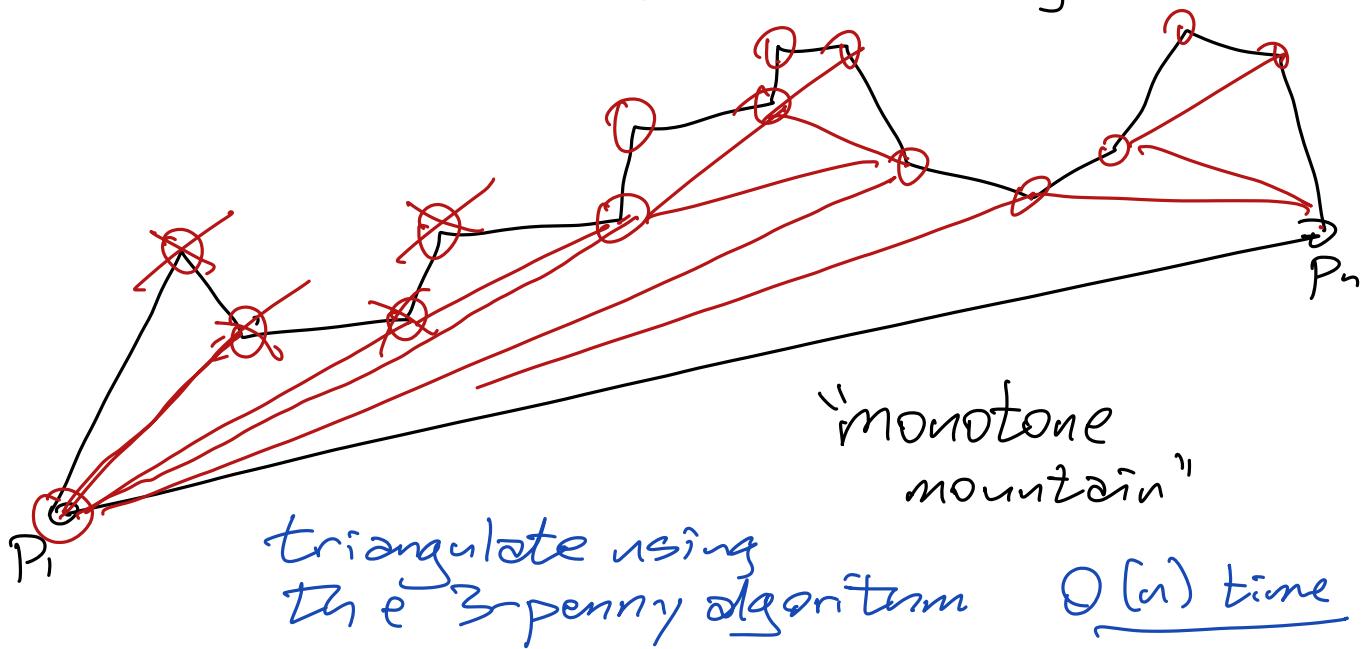
Trapezoid is boring unless
both on floor
or both on ceiling

Suppose no boring trapezoids

Every boring trap
contains a diagonal!



Linear sequence of interesting traps
wlog all with defining pts
on the ceiling



triangulate using
the 3-penny algorithm

$O(n)$ time

